

CLAIM AMENDMENTS:

1. (currently amended) A computer implemented method ~~Method~~ for optimizing the order of assignment of a number of supplies or resources to a number of demanders or demands, where each supply or resource has a certain supply or resource amount and each demander has a certain demand amount which is to be satisfied by said supplies or resources, the method comprising the steps of:

- building, by the computer, a network,

in which the supplies are represented by supply vertices connected to a sink vertex via sink edges of a flow capacity which represents the supply amount of the respective supply, the sink vertex being sink of a network flow;

in which the demanders are represented by demander vertices connected to a source vertex via source edges of a flow capacity which represents the demand amount of the respective demander, the source vertex being source of a network flow; and

in which supply vertices and demander vertices are connected by edges of certain flow capacities;

- determining, by the computer, an optimized network flow distribution of flow values through the edges by an iterative flow-method comprising the following steps a) to e):

- a) ~~repeated-repeating,~~ by the computer, a discharge operation pushing a flow for all active demander vertices, which are defined as vertices at which the sum of the incoming flow is higher than the sum of the outgoing flow along an edge being admissible according to a labeling

function, where the following rule applies: If the excess flow to be pushed from an active demander vertex is greater than the residual capacity of some current edge where the residual capacity of an edge is the capacity of the edge minus the flow that has already been pushed along this edge, and if there is another admissible edge leaving this vertex with a residual capacity which at least equals the excess flow, then the flow will be pushed along the edge with sufficient residual capacity;

- b) pushing, by the computer, excess flow from security vertices to the sink along all edges with a non zero residual capacity;
- c) setting, by the computer, all supply vertices with excess flow active and updating the excess flow of each vertex;
- d) ~~repeated-repeating, by the computer,~~ a discharge operation for all active supply vertices where flows are pushed back to demander vertices and where the following rules apply:
 - 1. a flow on an edge leaving a demander vertex that has no other edge leaving this demander vertex and leading to another supply vertex is never pushed back;
 - 2. if the flow that had been pushed along a current edge from the demander vertex is equal to the capacity of the edge then other edges that are admissible according to a labeling function would have priority under the condition that:

- i) If there is another admissible edge according to the labeling function from some demander vertex to the supply vertex with a flow that is less than the capacity of the edge and at least equal to the excess flow then the excess flow of supply vertex is pushed back along this edge to the demander vertex;
- ii) if there are other admissible edges to this supply vertex according to the labeling function where for each edge the flow is less than the capacity and where the sum of flows along these other edges is at least equal to the excess flow of the supply vertex then the excess flow is pushed back along these edges to the respective demander vertices;

e) setting, by the computer, demander vertices with excess flow after the reflow from active supply vertices ~~active in step d)~~ and updating the excess flow of each vertex;

- where the steps a) to e) are iteratively repeated until the flow to the sink vertex equals a min s-t-cut value or there is no more active vertex or the number of iterations has reached a maximum value;

and

- deriving, by the computer, the optimized order of assignment from the optimized network flow distribution by assigning the supply vertices to the demander vertices in correspondence to the flow values of the connecting edges.

Claims 2 - 7 (canceled).

8. (previously presented) Method according to claim 1, wherein the certain flow capacity of an edge connecting a demander vertex to a supply vertex is given by the smaller one of the capacity of the respective source edge and the capacity of the respective sink edge.

9. (original) Method according to claim 8, wherein the assigning operation, in a first stage, assigns a supply vertex to a demander vertex only if these vertices are connected by an edge for which the flow value equals the capacity.

10. (original) Method according to claim 9, wherein the assigning operation first assigns supply vertices to such demander vertices which are connected to the respective supply vertex by an edge for which the flow value equals the flow value of the corresponding source edge before it assigns supply vertices to such demander vertices which are connected to the respective supply vertex by an edge for which the flow value is equal to or higher than a remaining flow value of the corresponding sink edge which has not yet been assigned to a demander vertex.

11. (previously presented) Method according to claim 10, wherein the first stage is performed until all supply vertices and demander vertices which are connected by edges for which the flow value equals the capacity are assigned.

12. (original) Method according to claim 11, wherein the assigning operation, in a second stage, assigns a supply vertex to a demander vertex if the flow value of the connecting edge corresponds to the flow value of the corresponding source edge reduced by a fraction of its demand amount already assigned to a supply vertex, or

to the flow value of the corresponding sink edge reduced by a fraction of its supply amount already assigned to a demander vertex.

13. (original) Method according to claim 12, wherein the assigning operation, in the second stage, first assigns such supply vertices to demander vertices for which the flow value of the connecting edge corresponds to the flow value of the corresponding source edge reduced by a fraction of its demand amount which is already assigned to a supply vertex.

14. (previously presented) Method according to claim 1, wherein the certain flow capacity of an edge connecting a demander vertex to a supply vertex is given by the smaller one of the capacity of the respective source edge and the capacity of the respective sink edge.

15. (currently amended) A computer implemented method ~~Method~~ for balancing a number of loan accounts with a number of collateral securities, where each loan account has a certain loan value and each collateral security has a certain security value and wherein the collateral securities are to be offset against the loan accounts, the method comprising the steps of:

- building, by the computer, a network

in which the collateral securities are represented by security vertices connected to a sink vertex via sink edges of a flow capacity which represents the security value of the respective collateral security, the sink vertex being sink of a network flow;

in which the loan accounts are represented by loan vertices connected to a source vertex via source edges of a flow capacity which represents the loan value of the respective loan account, the source vertex being source of a network flow; and

in which security vertices and loan vertices are connected by edges of certain flow capacities;

- determining, by the computer, an optimized network flow distribution of flow values through the edges by an iterative flow-method comprising the following steps a) to e):

- a) ~~repeated-repeating~~, by the computer, a discharge operation pushing a flow for all active loan vertices, which are defined as vertices at which the sum of the incoming flow is higher than the sum of the outgoing flow along an edge being admissible according to a labeling function, where the following rule applies: If ~~the~~ excess flow to be pushed from an active loan vertex is greater than the residual capacity of some current edge where the residual capacity of an edge is the capacity of the edge minus the flow that has already been pushed along this edge, and if there is another admissible edge leaving this vertex with a residual capacity which at least equals the excess flow, then the flow will be pushed along the edge with sufficient residual capacity;
- b) pushing, by the computer, excess flow from security vertices to the sink along all edges with a non zero residual capacity;
- c) setting, by the computer, all security vertices with excess flow active and updating the excess flow of each vertex;

- d) repeated-repeating, by the computer, a discharge operation for all active security vertices where flows are pushed back to loan vertices and where the following rules apply:
1. a flow on an edge leaving a loan vertex that has no other edge leaving this loan vertex and leading to another security vertex is never pushed back;
 2. if the flow that had been pushed along a current edge from the loan vertex is equal to the capacity of the edge then other edges that are admissible according to a labeling function would have priority under the condition that:
 - i) If there is another admissible edge according to the labeling function from some loan vertex to the security vertex with a flow that is less than the capacity of the edge and at least equal to the excess flow then the excess flow of security vertex is pushed back along this edge to the loan vertex;
 - ii) if there are other admissible edges to this security vertex according to the labeling function where for each edge the flow is less than the capacity and where the sum of flows along these other edges is at least equal to the excess flow of the security vertex then the excess flow is pushed back along these edges to the respective loan vertices;

e) setting, by the computer, loan vertices with excess flow after the reflow from active security vertices-active in step d) and updating the excess flow of each vertex;

- where the steps a) to e) are iteratively repeated until the flow to the sink vertex equals a min s-t-cut value or there is no more active vertex or the number of iterations has reached a maximum value;

and

- deriving, by the computer, the optimized order of assignment from the optimized network flow distribution by offsetting the security vertices against the loan vertices in correspondence to the flow values of the connecting edges.

Claims 16 - 18 (canceled).

19. (previously presented) Method according to claim 15, comprising discharge operations pushing flows from loan vertices to security vertices and discharge operations pushing flows from security vertices to loan vertices.

20. (original) Method according to claim 19, wherein the discharge operation is performed iteratively for loan vertices and security vertices.

Claim 21 (canceled).

22. (previously presented) Method according to claim 15, wherein offsetting the security vertices against the loan vertices is performed by an iterative offsetting operation.

23. (original) Method according to claim 22, wherein the offsetting operation in a first stage offsets a security vertex against a loan vertex only if these vertices are connected by an edge for which the flow value equals the capacity.

24. (original) Method according to claim 23, wherein the offsetting operation first offsets security vertices against such loan vertices which are connected to the respective security vertex by an edge for which the flow value equals the flow value of the corresponding source edge before it offsets security vertices against such loan vertices which are connected to the respective security vertex by an edge for which the flow value is equal to or higher than a remaining flow value of the corresponding sink edge which has not yet been offset against a loan vertex.

25. (previously presented) Method according to claim 24, wherein the first stage is performed until all security vertices and loan vertices which are connected by edges for which the flow value equals the capacity are offset.

26. (original) Method according to claim 25, wherein the offsetting operation, in a second stage, offsets a security vertex against a loan vertex if the flow value of the edge connecting the respective vertices corresponds to the flow value of the corresponding source edge reduced by a fraction of its loan value already offset against a security vertex, or to the flow value of the corresponding sink edge reduced by a fraction of its security value already offset against a loan vertex.

27. (original) Method according to claim 26, wherein the offsetting operation, in the second stage, first offsets such security vertices against loan vertices for which the flow value of the edge connecting the respective vertices corresponds to the flow value of the corresponding source edge reduced by the loan value already offset against a collateral security.

28. (previously presented) Method according to claim 15, wherein the certain flow capacity of an edge connecting a loan vertex to a sink vertex is given by the

smaller one of the capacity of the respective source edge and the capacity of the respective sink edge.

29. (currently amended) A computer implemented method ~~Method~~ for optimizing the data transfer through a transmission system comprising a number of senders and a number of transmission lines, where each transmission line has a certain transmission rate and each sender is connected to a number of said transmission lines and has a certain data rate, by assigning transmission lines to senders, the method comprising the steps of:

- building, by the computer, a network

in which the transmission lines are represented by transmission vertices connected to a sink vertex via sink edges of a flow capacity which represents the transmission rate of the respective transmission line, the sink vertex being sink of a network flow;

in which the senders are represented by sender vertices connected to a source vertex via source edges of a flow capacity which represents the data rate of the respective sender, the source vertex being source of a network flow; and

in which transmission vertices and sender vertices are connected by edges of certain flow capacities;

- determining, by the computer, an optimized network flow distribution of flow values through the edges by an iterative flow-method comprising the following steps a) to e):

- a) repeated-repeating, by the computer, a discharge operation pushing a flow for all active sender vertices, which are defined as vertices at

which the sum of the incoming flow is higher than the sum of the outgoing flow along an edge being admissible according to a labeling function, where the following rule applies: If the excess flow to be pushed from an active sender vertex is greater than the residual capacity of some current edge where the residual capacity of an edge is the capacity of the edge minus the flow that has already been pushed along this edge, and if there is another admissible edge leaving this vertex with a residual capacity which at least equals the excess flow, then the flow will be pushed along the edge with sufficient residual capacity;

- b) pushing, by the computer, excess flow from transmission vertices to the sink along all edges with a non zero residual capacity;
- c) setting, by the computer, all transmission vertices with excess flow active and updating the excess flow of each vertex;
- d) repeated-repeating, by the computer, a discharge operation for all active transmission vertices where flows are pushed back to sender vertices and where the following rules apply:
 - 1. a flow on an edge leaving a sender vertex that has no other edge leaving this sender vertex and leading to another transmission vertex is never pushed back;
 - 2. if the flow that had been pushed along a current edge from the sender vertex is equal to the capacity of the edge then other edges that are admissible according to a labeling function

would have priority under the condition that:

- i) If there is another admissible edge according to the labeling function from some sender vertex to the transmission vertex with a flow that is less than the capacity of the edge and at least equal to the excess flow then the excess flow of transmission vertex is pushed back along this edge to the sender vertex;
- ii) if there are other admissible edges to this transmission vertex according to the labeling function where for each edge the flow is less than the capacity and where the sum of flows along these other edges is at least equal to the excess flow of the transmission vertex then the excess flow is pushed back along these edges to the respective sender vertices;

e) setting, by the computer, sender vertices with excess flow after the reflow from active transmission vertices active-in step d) and updating the excess flow of each vertex;

- where the steps a) to e) are iteratively repeated until the flow the to sink vertex equals a min s-t-cut value or there is no more active vertex or the number of iterations has reached a maximum value;

and

- deriving, by the computer, the optimized order of assignment from the optimized network flow distribution by assigning the transmission vertices to the sender vertices in correspondence to the flow values of the connecting edges.

30. (original) Method according to claim 29, wherein in the iterative flow-method comprises a discharge operation pushing a flow from an active vertex at which the sum of the incoming network flow is higher than the sum of the outgoing network flow along an admissible edge, where the admissibility of an edge is defined by a label of the vertex connected to the active vertex by the respective edge.

31. (original) Method according to claim 30, further comprising a relabeling operation changing the label of the active vertex if there is no admissible edge along which the discharge operation can be performed.

32. (original) Method according to claim 31, wherein, when the label of the vertex to be discharged is $\Psi(v)$ and the label of a vertex connected by an edge is $\Psi(w)$, said edge being admissible if $\Psi(v) = \Psi(w) + 1$, and wherein the label $\Psi(v)$ of the vertex to be discharged is increased by one in the relabeling operation.

33. (previously presented) Method according to claim 32, comprising discharge operations pushing flows from sender vertices to transmission vertices and discharge operations pushing flows from transmission vertices to sender vertices.

34. (original) Method according to claim 33, wherein the discharge operation is performed iteratively for sender vertices and transmission vertices.

35. (previously presented) Method according to claim 29, wherein a total flow is the network flow through the edges from the source vertex to the sink vertex and the flow-method comprises the steps of:

- determining an upper limit of the highest possible total flow through the edges; and

- iteratively distributing the network flow through the edges until at least one of the conditions is fulfilled:

i) the network flow corresponds to the upper limit of the highest possible total flow,

ii) the sum of the incoming network flow at a vertex equals the sum of the outgoing network flow of said vertex for each transmission vertex and for each sender vertex,

iii) the number of iterations has reached a given maximum value.

36. (previously presented) Method according to claim 29, wherein assigning the transmission vertices to the sender vertices is performed by an iterative assigning operation.

37. (original) Method according to claim 36, wherein the assigning operation, in a first stage, assigns a transmission vertex to a sender vertex only if these vertices are connected by an edge for which the flow value equals the capacity.

38. (original) Method according to claim 37, wherein the assigning operation first assigns transmission vertices to such sender vertices which are connected to the respective transmission vertex by an edge for which the flow value equals the flow value of the corresponding source edge before it assigns transmission vertices to such sender vertices which are connected to the respective supply vertex by an edge for which the flow value is equal to or higher than a remaining flow value of the corresponding sink edge which has not yet been assigned to a sender vertex.

39. (previously presented) Method according to claim 38, wherein the first stage is performed until all transmission vertices and sender vertices which are connected by edges for which the flow value equals the capacity are assigned.

40. (original) Method according to claim 39, wherein the assigning operation, in a second stage, assigns a transmission vertex to a sender vertex if the flow value of the connecting edge corresponds to the flow value of the corresponding source edge reduced by a fraction of its data rate already assigned to a transmission vertex, or to the flow value of the corresponding sink edge reduced by a fraction of its transmission rate already assigned to a sender vertex.

41. (original) Method according to claim 40, wherein the assigning operation, in the second stage, first assigns such transmission vertices to sender vertices for which the flow value of the connecting edge corresponds to the flow value of the corresponding source edge reduced by a fraction of its data rate already assigned to a transmission vertex.

42. (previously presented) Method according to claim 41, wherein the certain flow capacity of an edge connecting a sender vertex to a transmission vertex is given by the smaller one of the capacity of the respective source edge and the capacity of the respective sink edge.

43. (currently amended) A computer implemented method Method for optimizing the order of assignment of a number of tasks to a number of processors, where each processor has a certain processor capacity and each task has a certain capacity demand which is to be satisfied by at least one of said processors, the method comprising the steps of:

- building, by the computer, a network

in which the processors are represented by processor vertices connected to a sink vertex via sink edges of a flow capacity which represents the processor capacity of the respective processor, the sink vertex being sink of a network flow;

in which the tasks are represented by task vertices connected to a source vertex via source edges of a flow capacity which represents the capacity demand of the respective task, the source vertex being source of a network flow; and

in which processor vertices and task vertices are connected by edges of certain flow capacities;

- determining, by the computer, an optimized network flow distribution of flow values through the edges by an iterative flow-method comprising the following steps a) to e):

- a) repeated-repeating, by the computer, a discharge operation pushing a flow for all active task vertices, which are defined as vertices at which the sum of the incoming flow is higher than the sum of the outgoing flow along an edge being admissible according to a labeling function, where the following rule applies: If the excess flow to be pushed from an active task vertex is greater than the residual capacity of some current edge where the residual capacity of an edge is the capacity of the edge minus the flow that has already been pushed along this edge, and if there is another admissible edge leaving this vertex with a residual capacity which at least equals the excess flow,

then the flow will be pushed along the edge with sufficient residual capacity;

- b) pushing, by the computer, excess flow from processor vertices to the sink along all edges with a non zero residual capacity;
- c) setting, by the computer, all processor vertices with excess flow active and updating the excess flow of each vertex;
- d) repeated-repeating, by the computer, a discharge operation for all active processor vertices where flows are pushed back to task vertices and where the following rules apply:
 - 1. a flow on an edge leaving a task vertex that has no other edge leaving this task vertex and leading to another processor vertex is never pushed back;
 - 2. if the flow that had been pushed along a current edge from the task vertex is equal to the capacity of the edge then other edges that are admissible according to a labeling function would have priority under the condition that:
 - i) If there is another admissible edge according to the labeling function from some task vertex to the processor vertex with a flow that is less than the capacity of the edge and at least equal to the excess flow then the excess flow of processor vertex is pushed back along this edge to the task vertex;

- ii) if there are other admissible edges to this processor vertex according to the labeling function where for each edge the flow is less than the capacity and where the sum of flows along these other edges is at least equal to the excess flow of the processor vertex then the excess flow is pushed back along these edges to the respective task vertices;

e) setting, by the computer, task vertices with excess flow after the reflow from active processor vertices active in step d) and updating the excess flow of each vertex;

- where the steps a) to e) are iteratively repeated until the flow to the sink vertex equals a min s-t-cut value or there is no more active vertex or the number of iterations has reached a maximum value;

and

- deriving, by the computer, the optimized order of assignment from the optimized network flow distribution by assigning the processor vertices to the task vertices in correspondence to the flow values of the connecting edges.

44. (original) Method according to claim 43, wherein in the iterative flow-method comprises a discharge operation pushing a flow from an active vertex at which the sum of the incoming network flow is higher than the sum of the outgoing network flow along an admissible edge, where the admissibility of an edge is defined by a label of the vertex connected to the active vertex by the respective edge.

45. (original) Method according to claim 44, further comprising a relabeling operation changing the label of the active vertex if there is no admissible edge along which the discharge operation can be performed.

46. (previously presented) Method according to claim 45, wherein, when the label of the vertex to be discharged is $\Psi(v)$ and the label of a vertex connected by an edge is $\Psi(w)$, said edge being admissible if $\Psi(v) = \Psi(w) + 1$, and wherein the label $\Psi(v)$ of the vertex to be discharged is increased by one in the relabeling operation.

47. (previously presented) Method according to claim 46, comprising discharge operations pushing flows from task vertices to processor vertices and discharge operations pushing flows from processor vertices to task vertices.

48. (original) Method according to claim 47, wherein the discharge operation is performed iteratively for task vertices and processor vertices.

49. (previously presented) Method according to claim 43, wherein a total flow is the network flow through the edges from the source vertex to the sink vertex and the flow-method comprises the steps of:

- determining an upper limit of the highest possible total flow through the edges; and
- iteratively distributing the network flow through the edges until at least one of the conditions is fulfilled:
 - i) the network flow corresponds to the upper limit of the highest possible total flow,

ii) the sum of the incoming network flow at a vertex equals the sum of the outgoing network flow of said vertex for each processor vertex and for each task vertex,

iii) the number of iterations has reached a given maximum value.

50. (previously presented) Method according to claim 49, wherein assigning the processor vertices to the task vertices is performed by an iterative assigning operation.

51. (original) Method according to claim 50, wherein the assigning operation, in a first stage, assigns a processor vertex to a task vertex only if these vertices are connected by an edge for which the flow value equals the capacity.

52. (original) Method according to claim 51, wherein the assigning operation first assigns processor vertices to such task vertices which are connected to the respective processor vertex by an edge for which the flow value equals the flow value of the corresponding source edge before it assigns processor vertices to such task vertices which are connected to the respective processor vertex by an edge for which the flow value is equal to or higher than a remaining flow value of the corresponding sink edge which has not yet been assigned to a task vertex.

53. (previously presented) Method according to claim 52, wherein the first stage is performed until all processor vertices and task vertices which are connected by edges for which the flow value equals the capacity are assigned.

54. (original) Method according to claim 53, wherein the assigning operation, in a second stage, assigns a processor vertex to a task vertex if the flow value of the connecting edge corresponds to the flow value of the corresponding source edge

reduced by a fraction of its capacity demand already assigned to a processor vertex, or to the flow value of the corresponding sink edge reduced by a fraction of its processing capacity already assigned to a task vertex.

55. (original) Method according to claim 54, wherein the assigning operation in the second stage first assigns such processor vertices to task vertices for which the flow value of the connecting edge corresponds to the flow value of the corresponding source edge reduced by a fraction of its capacity demand which is already assigned to a processor vertex.

56. (previously presented) Method according to claim 43, wherein the certain flow capacity of an edge connecting a task vertex to a processor vertex is given by the smaller one of the capacity of the respective source edge and the capacity of the respective sink edge.

57. (currently amended) ~~Device~~ A computer device for determining an optimized assignment of a number of supplies or resources, each having a certain supply or resource amount to a number of demanders or demands, each having a certain demand amount to be satisfied by said supplies, in which, after the assignment, the sum of unsatisfied demand amounts is minimized, comprising:

- a supply input computer unit for inputting supply data representing supplies and their supply amounts;
- a demander input computer unit for inputting demander data representing demanders and their demand amounts,

- an access input computer unit for inputting access data representing, for each demander, the corresponding supplies which can be accessed by the respective demander for satisfying its demand amount;

- a network construction computer unit for constructing, on the basis of the supply data, the demander data and the access data, a network comprising:

- a) a supply vertex for each supply,
- b) a demander vertex for each demander,
- c) a sink vertex,
- d) a source vertex,
- e) edges, each having a certain flow capacity and connecting a supply vertex and a demander vertex,
- f) sink edges, each connecting the sink vertex to one of the supply vertices and having a flow capacity representing the supply amount of the respective supply, and
- g) source edges, each connecting the source vertex to one of the demander vertices and having a flow capacity representing the demand amount of the respective demander;

- a network flow computer unit for determining an optimized network flow distribution through the network, the optimized network flow being represented by flow values through the edges, the network flow being designed to determine an optimized network flow distribution of flow values through the edges further comprising:

- a) means for ~~repeated-repeating, by the computer device, a~~ discharge operation pushing a flow for all active demander vertices, which are

defined as vertices at which the sum of the incoming flow is higher than the sum of the outgoing flow along an edge being admissible according to a labeling function, where the following rule applies: If the excess flow to be pushed from an active demander vertex is greater than the residual capacity of some current edge where the residual capacity of an edge is the capacity of the edge minus the flow that has already been pushed along this edge, and if there is another admissible edge leaving this vertex with a residual capacity which at least equals the excess flow, then the flow will be pushed along the edge with sufficient residual capacity;

- b) means for pushing, by the computer device, excess flow from security vertices to the sink along all edges with a non zero residual capacity;
- c) means for setting, by the computer device, all supply vertices with excess flow active and updating the excess flow of each vertex;
- d) means for ~~repeated-repeating~~, by the computer device, discharge operation for all active supply vertices where flows are pushed back to demander vertices and where the following rules apply:
 - 1. a flow on an edge leaving a demander vertex that has no other edge leaving this demander vertex and leading to another supply vertex is never pushed back;
 - 2. if the flow that had been pushed along a current edge from the demander vertex is equal to the capacity of the edge then other edges that are admissible according to a labeling function

would have priority under the condition that:

- i) If there is another admissible edge according to the labeling function from some demander vertex to the supply vertex with a flow that is less than the capacity of the edge and at least equal to the excess flow then the excess flow of supply vertex is pushed back along this edge to the demander vertex;
- ii) if there are other admissible edges to this supply vertex according to the labeling function where for each edge the flow is less than the capacity and where the sum of flows along these other edges is at least equal to the excess flow of the supply vertex then the excess flow is pushed back along these edges to the respective demander vertices;

e) means for setting, by the computer device, demander vertices with excess flow after the reflow from active supply vertices active by means d) and updating the excess flow of each vertex;

- where the means a) to e) are iteratively repeated until the flow to the sink vertex equals a min s-t-cut value or there is no more active vertex or the number of iterations has reached a maximum value;

and

- an assignment computer unit for assigning the supplies to the demanders by assigning the supply vertices to the demander vertices in correspondence to the flow values of the connecting edges.

58. (currently amended) Device according to claim 57, wherein the input computer units are formed by a single input computer unit.

59. (currently amended) Device according to claim 57, wherein the input computer units are integrated into a single device.

60. (currently amended) Device according to claim 59, wherein the network construction computer unit, the network flow computer unit, and the assignment computer unit are realized by a single calculator computer unit.

61. (currently amended) ~~Device~~—A computer device for balancing a number of loan accounts with a number of collateral securities, where each loan account has a certain loan value and each collateral security has a certain security value and wherein the collateral securities are to be offset against the loan accounts, comprising:

- a security input computer unit for inputting security data representing collateral securities and their security values;

- a loan input computer unit for inputting loan data representing loan accounts and their loan values,

- an access input computer unit for inputting access data representing, for each loan account, the corresponding collateral securities, which can be offset against the respective loan account;

- a network construction computer unit for constructing, on the basis of the security data, the loan data and the access data, a network comprising:

- a) a security vertex for each security,
- b) a loan vertex for each loan,
- c) a sink vertex,
- d) a source vertex,
- e) edges, each having a certain flow capacity and connecting a security vertex and a loan vertex,

- f) sink edges, each connecting the sink vertex to one of the security vertices and having a flow capacity representing the security value of the respective collateral security, and

- g) source edges, each connecting the source vertex to one of the loan vertices and having a flow capacity representing the loan value of the respective loan account;

- a network flow unit for determining an optimized network flow distribution through the network, the optimized network flow being represented by flow values through the edges, the network flow being designed to determine an optimized network flow distribution of flow values through the edges further comprising:

- a) means for ~~repeated-repeating, by the computer device, a~~ discharge operation pushing a flow for all active loan vertices, which are defined as vertices at which the sum of the incoming flow is higher than the sum of the outgoing flow along an edge being admissible according to a labeling function, where the following rule applies: If ~~the~~ excess flow to be pushed from an active loan vertex is greater than the residual capacity of some current edge where the residual capacity of an edge

is the capacity of the edge minus the flow that has already been pushed along this edge, and if there is another admissible edge leaving this vertex with a residual capacity which at least equals the excess flow, then the flow will be pushed along the edge with sufficient residual capacity;

- b) means for pushing, by the computer device, excess flow from security vertices to the sink along all edges with a non zero residual capacity;
- c) means for setting, by the computer device, all security vertices with excess flow active and updating the excess flow of each vertex;
- d) means for ~~repeated~~-repeating, by the computer device, discharge operation for all active security vertices where flows are pushed back to loan vertices and where the following rules apply:
 - 1. a flow on an edge leaving a loan vertex that has no other edge leaving this loan vertex and leading to another security vertex is never pushed back;
 - 2. if the flow that had been pushed along a current edge from the loan vertex is equal to the capacity of the edge then other edges that are admissible according to a labeling function would have priority under the condition that:
 - i) If there is another admissible edge according to the labeling function from some loan vertex to the security vertex with a flow that is less than the capacity

of the edge and at least equal to the excess flow then the excess flow of security vertex is pushed back along this edge to the loan vertex;

- ii) if there are other admissible edges to this security vertex according to the labeling function where for each edge the flow is less than the capacity and where the sum of flows along these other edges is at least equal to the excess flow of the security vertex then the excess flow is pushed back along these edges to the respective loan vertices;

e) means for setting, by the computer device, loan vertices with excess flow after ~~the~~ reflow from active security vertices ~~active by means d)~~ and updating the excess flow of each vertex;

- where the means a) to e) are iteratively repeated until the flow to the sink vertex equals a min s-t-cut value or there is no more active vertex or the number of iterations has reached a maximum value;
and

- an assignment computer unit for offsetting the collateral securities against the loan accounts by assigning the security vertices to the loan vertices in correspondence to the flow values of the connecting edges.

62. (currently amended) Device according to claim 61, wherein the input computer units are formed by a single input computer unit.

63. (currently amended) Device according to claim 61, wherein the input

computer units are integrated into a single device.

64. (currently amended) Device according to claim 61, wherein the

network construction computer unit, the network flow computer unit, and the assignment computer unit are realized by a single calculator computer unit.

65. (currently amended) ~~Device~~ A computer device for determining an

optimized data transfer through a transmission system comprising a number of senders and a number of transmission lines by assigning transmission lines to senders, where each transmission line has a certain transmission rate and each sender is connected to a number of said transmission lines and has a certain data rate, comprising:

- a transmission line input computer unit for inputting transmission line data representing transmission lines and their transmission rates,

- a sender input computer unit for inputting sender data representing senders and their data rates,

- an access input computer unit for inputting access data representing, for each sender, the corresponding transmission lines, which can be accessed by a sender for transferring data;

- a network construction computer unit for constructing, on the basis of the transmission line data, the sender data and the access data, a network comprising:

- a) a transmission vertex for each transmission line,
- b) a sender vertex for each sender,
- c) a sink vertex,
- d) a source vertex,

e) edges, each having a certain flow capacity and connecting a transmission vertex and a sender vertex,

f) sink edges, each connecting the sink vertex to one of the transmission vertices and having a flow capacity representing the transmission rate of the respective transmission line, and

g) source edges, each connecting the source vertex to one of the sender vertices and having a flow capacity representing the data rate of the respective sender;

- a network flow computer unit for determining an optimized network flow distribution through the network, the optimized network flow being represented by flow values through the edges, the network flow being designed to determine an optimized network flow distribution of flow values through the edges further comprising:

a) means for ~~repeated~~ repeating, by the computer device, a discharge operation pushing a flow for all active sender vertices, which are defined as vertices at which the sum of the incoming flow is higher than the sum of the outgoing flow along an edge being admissible according to a labeling function, where the following rule applies: If the excess flow to be pushed from an active sender vertex is greater than the residual capacity of some current edge where the residual capacity of an edge is the capacity of the edge minus the flow that has already been pushed along this edge, and if there is another admissible edge leaving this vertex with a residual capacity which at least equals the excess flow, then the flow will be pushed along the edge with sufficient residual capacity;

- b) means for pushing, by the computer device, excess flow from transmission vertices to the sink along all edges with a non zero residual capacity;
- c) means for setting, by the computer device, all transmission vertices with excess flow active and updating the excess flow of each vertex;
- d) means for ~~repeated~~ repeating, by the computer device, discharge operation for all active transmission vertices where flows are pushed back to sender vertices and where the following rules apply:
 1. a flow on an edge leaving a sender vertex that has no other edge leaving this sender vertex and leading to another transmission vertex is never pushed back;
 2. if the flow that had been pushed along a current edge from the sender vertex is equal to the capacity of the edge then other edges that are admissible according to a labeling function would have priority under the condition that:
 - i) If there is another admissible edge according to the labeling function from some sender vertex to the transmission vertex with a flow that is less than the capacity of the edge and at least equal to the excess flow then the excess flow of transmission vertex is pushed back along this edge to the sender vertex;
 - ii) if there are other admissible edges to this transmission vertex according to the labeling function where for each

edge the flow is less than the capacity and where the sum of flows along these other edges is at least equal to the excess flow of the transmission vertex then the excess flow is pushed back along these edges to the respective sender vertices;

e) means for setting, by the computer device, sender vertices with excess flow after the reflow from active transmission vertices active by means d) and updating the excess flow of each vertex;

- where the means a) to e) are iteratively repeated until the flow to the sink vertex equals a min s-t-cut value or there is no more active vertex or the number of iterations has reached a maximum value;

and

- an assignment computer unit for assigning the transmission lines to the senders by assigning the transmission vertices to the sender vertices in correspondence to the flow values of the connecting edges.

66. (currently amended) Device according to claim 65, wherein the input computer units are formed by a single input computer unit.

67. (currently amended) Device according to claim 65, wherein the input computer units are integrated into a single device.

68. (currently amended) Device according to claim 65, wherein the network construction computer unit, the network flow computer unit, and the assignment computer unit are realized by a single calculator computer unit.

Claims 69 - 72 (canceled).

73. (previously presented) Computer program product for optimizing the order of assignment of a number of supplies to a number of demanders comprising instructions which, when loaded into a computer, cause said computer to perform a method as claimed in claim 1.

74. (previously presented) Computer program product for balancing a number of loan accounts with a number of collateral securities comprising instructions which, when loaded into a computer, cause said computer to perform a method as claimed in claim 15.

75. (previously presented) Computer program product for optimizing the data transfer through a transmission system comprising a number of senders and a number of transmission lines comprising instructions which, when loaded into a computer, cause said computer to perform a method as claimed in claim 29.

76. (previously presented) Computer program product for optimizing the order of assignment of a number of tasks to a number of processors comprising instructions which, when loaded into a computer, cause said computer to perform the method as claimed in claim 43.

77. (original) Storage medium comprising stored data which represent a computer program product as claimed in claim 73.

78. (original) Storage medium comprising stored data which represent a computer program product as claimed in claim 74.

79. (original) Storage medium comprising stored data which represent a computer program product as claimed in claim 75.

80. (original) Storage medium comprising stored data which represent a computer program product as claimed in claim 76.

81. (previously presented) Method according to claim 1, wherein the labeling function comprises the steps of relabeling an active vertex if during a discharge operation the excess flow of a vertex cannot be pushed because there are no admissible edges.

82. (previously presented) Method according to claim 81, wherein, in case of a flow, the label value of a demand vertex is updated to one plus the smallest label value of all supply vertices, and, in case of a reflow, the label value of a supply vertex is updated to one plus the smallest label value of all demand vertices.